

Concurrent Programming from pseuCo to Petri

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The Problem:



The Problem: Teaching Concurrency to Students



Ye Olde Concurrent Programming Lecture

— est. 2005 —

- 1. LTS: Labeled Transition Systems**
- 2. CCS: Calculus of Communicating Systems**
- 3. Concurrency in Java**

The Concurrent Programming Lecture before 2014



The Concurrent Programming Lecture since 2014



pseuCo: Bridging the Gap from Theory to Practical Programming since 2014

- minimal programming language
- Java-inspired Syntax
- both shared memory and message passing

```
void factorial(intchan c) {
    int z, j, n;
    while (true) {
        z = <? c; // receive input

        n = 1;
        for (j = z; j > 0 ; j--) {
            n = n*j;
        }

        c <! n; // send result
    };
}
```


pseuCo by Example

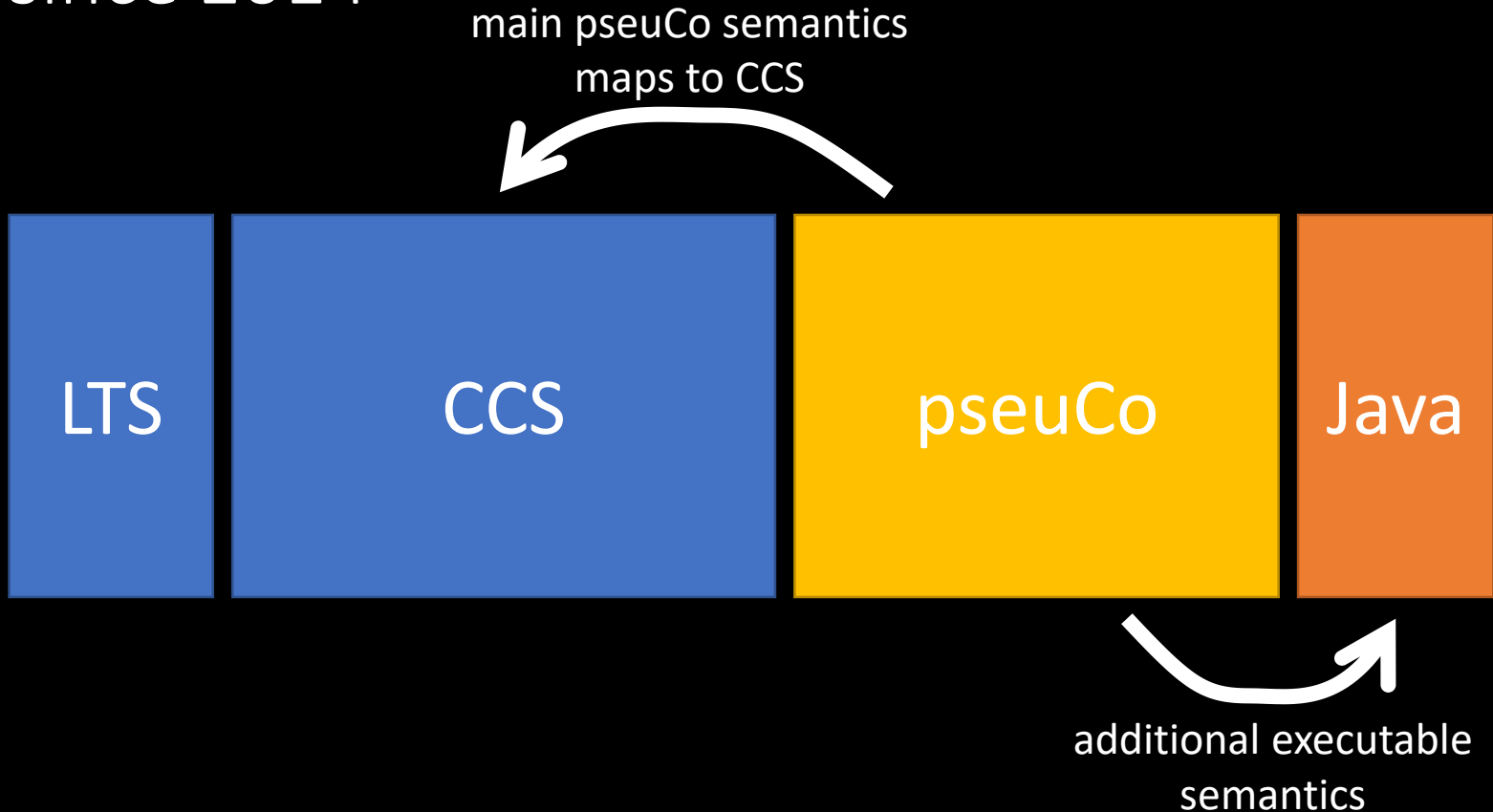
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    int z, j, n;
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        c <! n; // send result
    };
}
```

```
mainAgent {
    intchan cc;
    agent a = start(factorial(cc));
    cc <! 3;
    int mid = <? cc;
    println("3! evaluates to " + mid + ".");
    cc <! mid;
    println("(3!)! evaluates to " + (<? cc) + ".");
}
```

The Concurrent Programming Lecture since 2014



Editing

First Message Passing

pseuCo

Duplicate

pseuCo → CCS → LTS

Actions

pseuCo

1

No issues.

CCS

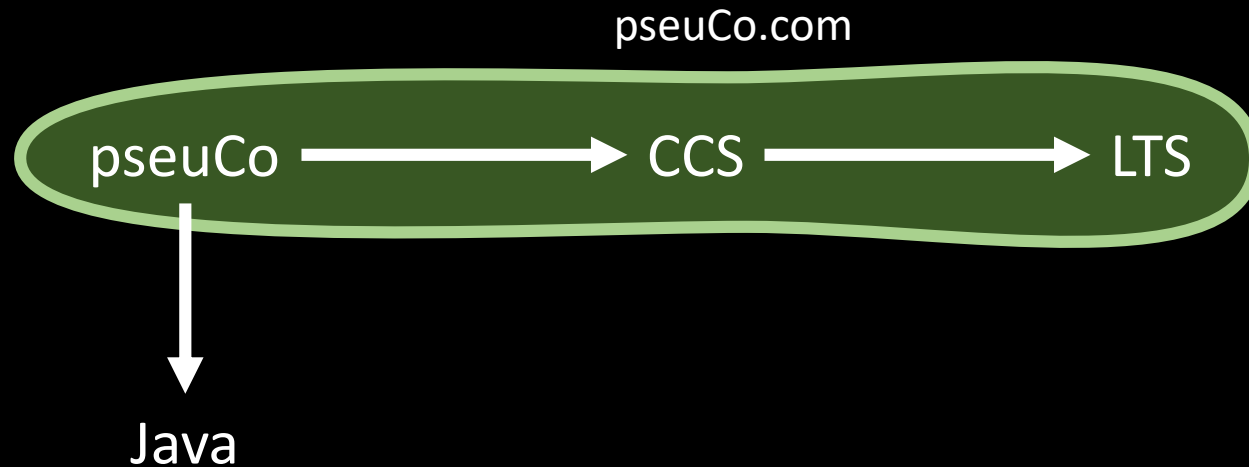
1

No issues.

LTS

+ Expand all

Where We Are



Contribution of this Paper: Beyond the pseuCo \rightarrow CCS Semantics

The pseuCo \rightarrow CCS translation has served us well, but it has problems:

pseuCo → CCS by example

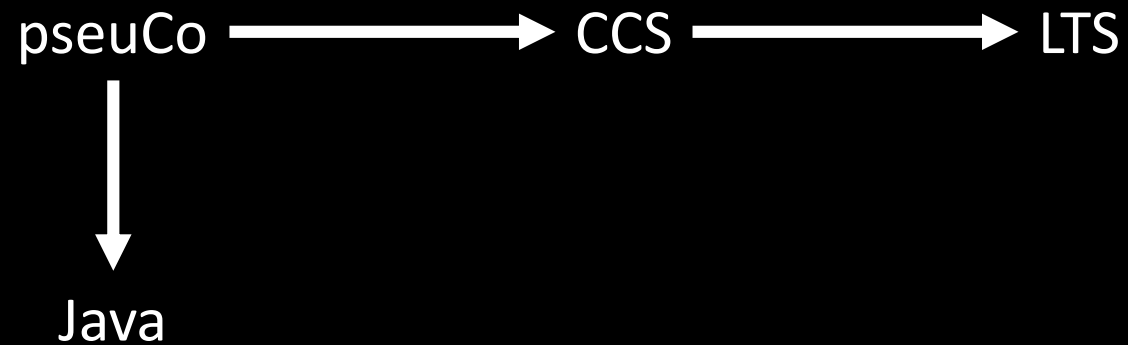
```
Channel_cons[i] := channel_create!(i).(Channel_cons[i-1])
AgentManager[next_i] := agent_new!(next_i).(AgentManager[next_i+1])
Agent_factorial :=
agent_new?i.(start_factorial!(i).(start_set_arg(i)?starter.(start_set_arg(i)?a1.(Agent_factorial |
(Proc_factorial[i, a1] ; agent_terminate(i)!.(0))))))
Proc_factorial[a, $c] := Proc_factorial_1[a, $c, 0, 0, 0]
Proc_factorial_1[a, $c, $z, $j, $n] := when (true) (τ.(receive($c)?$0.(Proc_factorial_3[a, $c, $0, $0, 1]))) +
when (!true) (Proc_factorial_6[a, $c, $z, $j, $n])
Proc_factorial_2[a, $c, $z, $j, $n] := Proc_factorial_3[a, $c, $z, $j-1, $n]
Proc_factorial_3[a, $c, $z, $j, $n] := when (!( $j > 0 )) (Proc_factorial_4[a, $c, $z, $j, $n]) + when ( $j > 0 )
(τ.(Proc_factorial_2[a, $c, $z, $j, $n * $j]))
Proc_factorial_4[a, $c, $z, $j, $n] := when ( $c >= 0 ) (put($c)!($n).(Proc_factorial_5[a, $c, $z, $j, $n])) + when
( $c < 0 ) (receive($c)!($n).(Proc_factorial_5[a, $c, $z, $j, $n]))
Proc_factorial_5[a, $c, $z, $j, $n] := Proc_factorial_1[a, $c, $z, $j, $n]
Proc_factorial_6[a, $c, $z, $j, $n] := 1
MainAgent[a] := channel_create?$0 (start_factorial?$1.(start_set_arg($1)!(a).(start_set_arg($1)!($0).(when
($0 >= 0) (put($0)!(3).(MainAgent_1[a, $0, $1])) + when ($0 < 0) (receive($0)!(3).(MainAgent_1[a, $0, $1]))))))
MainAgent_1[a, $cc, $a] := receive($cc)?$0 (println!("3! evaluates to " ^ $0 ^ " ").(when ($cc >= 0)
(put($cc)!($0).(MainAgent_2[a, $cc, $a, $0])) + when ($cc < 0) (receive($cc)!($0).(MainAgent_2[a, $cc, $a, $0]))))
MainAgent_2[a, $cc, $a, $mid] := receive($cc)?$0.(println!("(3!)! evaluates to " ^ $0 ^ " ").(0))
(Agent_factorial | MainAgent[1] | Channel_cons[-1] | AgentManager[2]) \ {*, println, exception}
```

Contribution of this Paper: Fixing the pseuCo \rightarrow CCS semantics

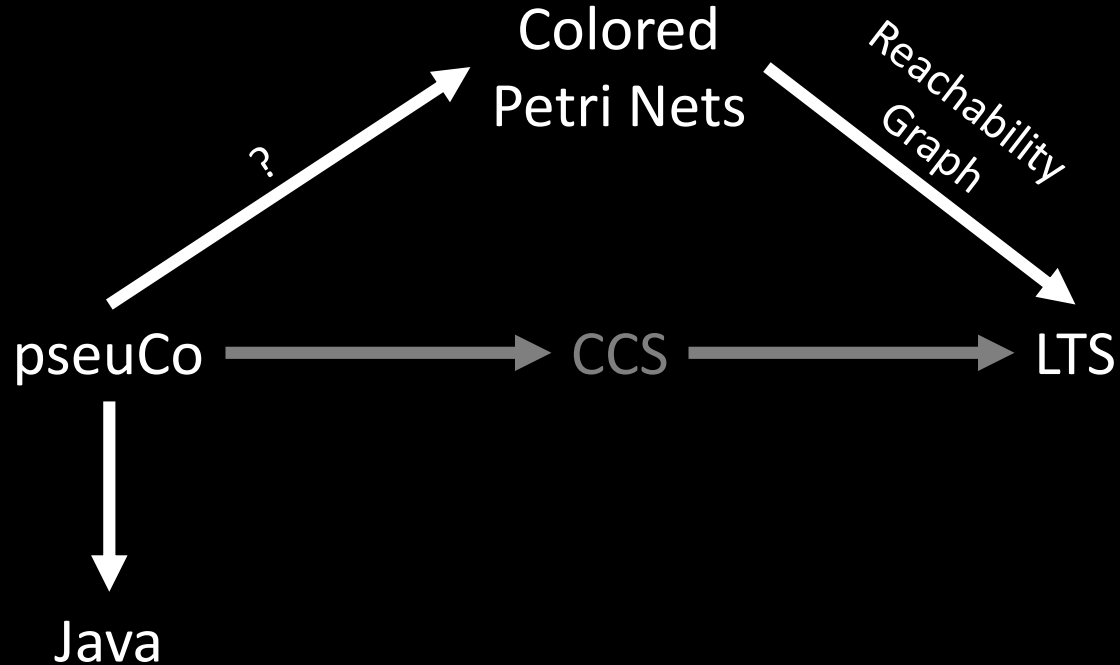
The pseuCo \rightarrow CCS translation has served us well, but it has problems:

- hard to understand
 - control flow hard to see (goto-style spaghetti code)
 - helper constructs (agent management, channels, arrays, ...)
 - low-level hacks visible in the code (e.g. for channels)
- no proper debugging support in pseuCo.com
- no true concurrency due to CCS interleaving semantics
- ...and lack of Petri Nets!

Where We Are



Where We Want to Go

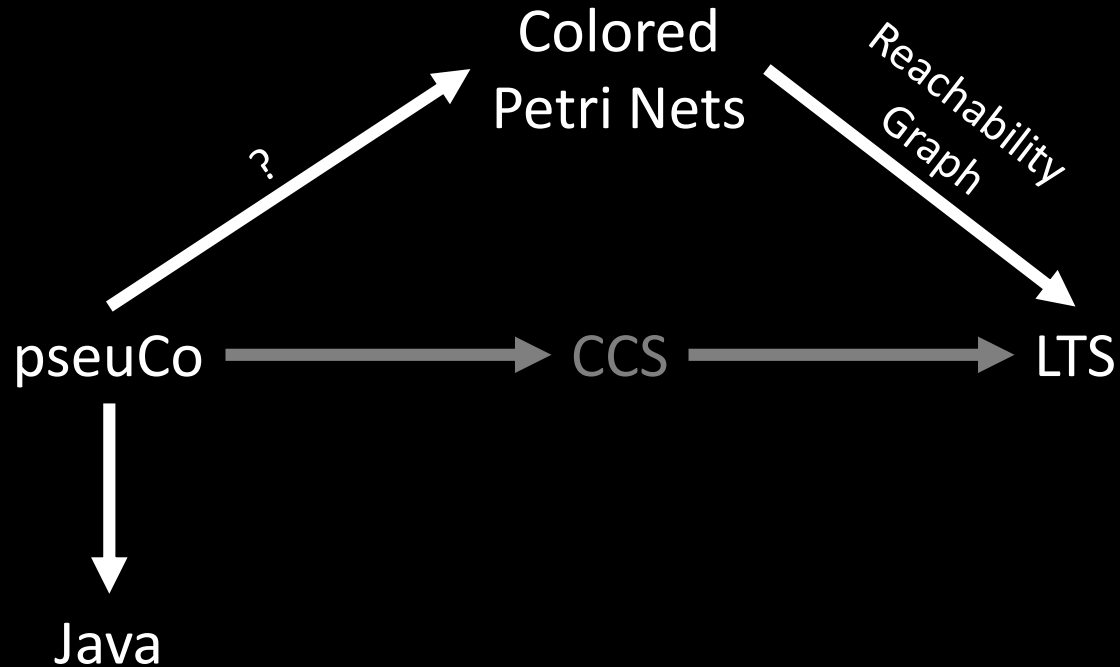


Petri Nets to the Rescue

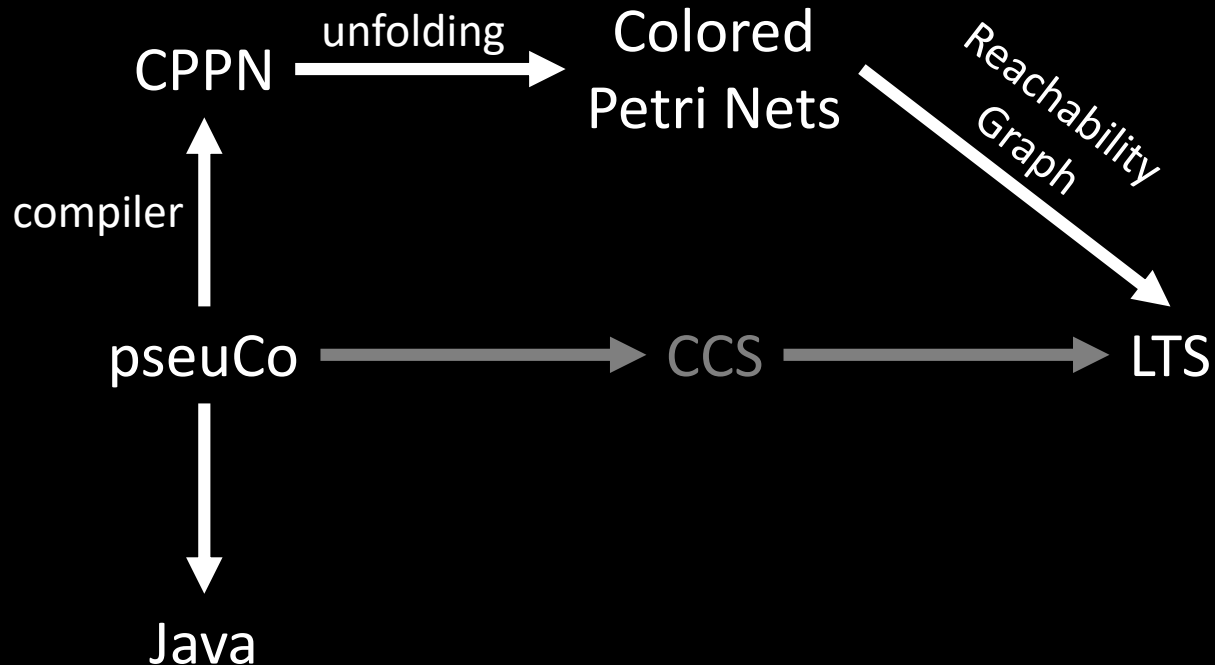
The pseuCo \rightarrow CCS translation has served us well, but it has problems:

- hard to understand
 - control flow hard to see (goto-style spaghetti code) ✓
 - helper constructs (agent management, channels, arrays, ...) ←
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- no true concurrency due to CCS interleaving semantics ✓
- ...and lack of Petri Nets! ✓

Where We Want to Go



Introducing Colored Program Petri Nets

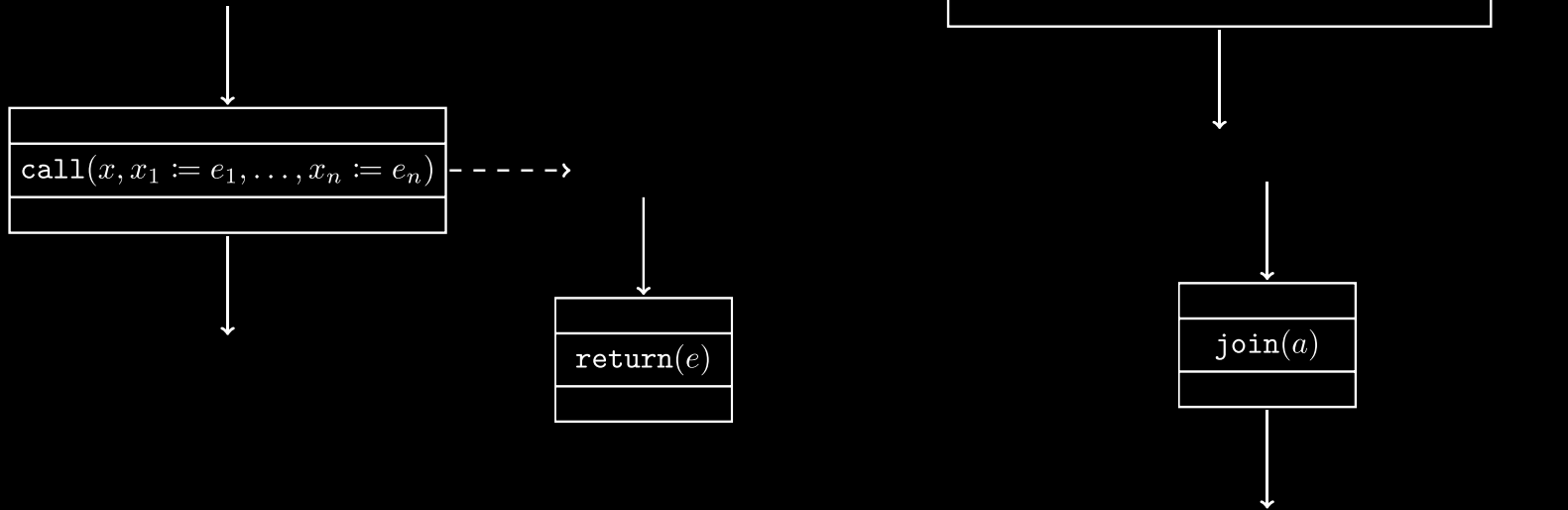


Introducing CPPN: Syntax

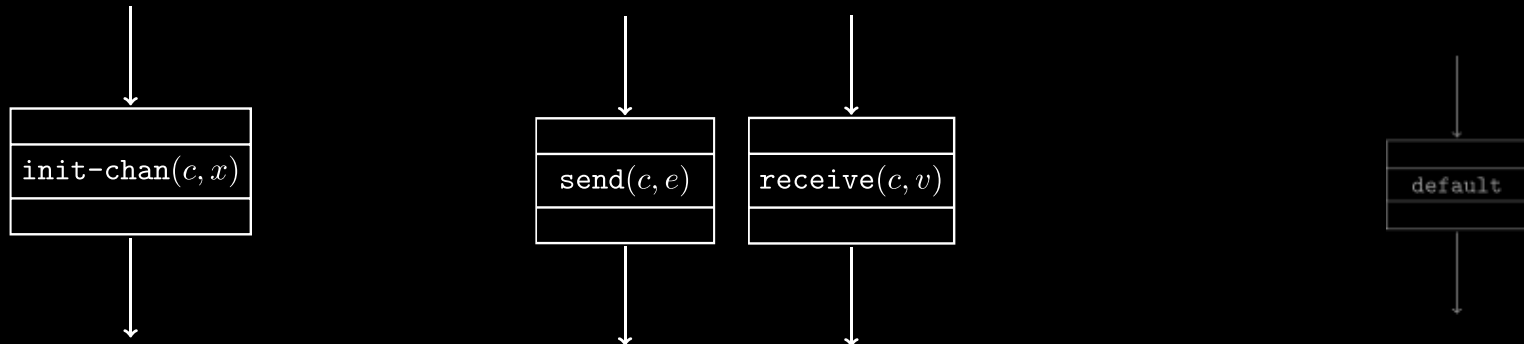
Syntax similar to CPN, but with two kinds of transitions:

- normal transitions (“Petri transitions”)
- *command* transitions
 - annotated with a command
 - fixed number of arcs

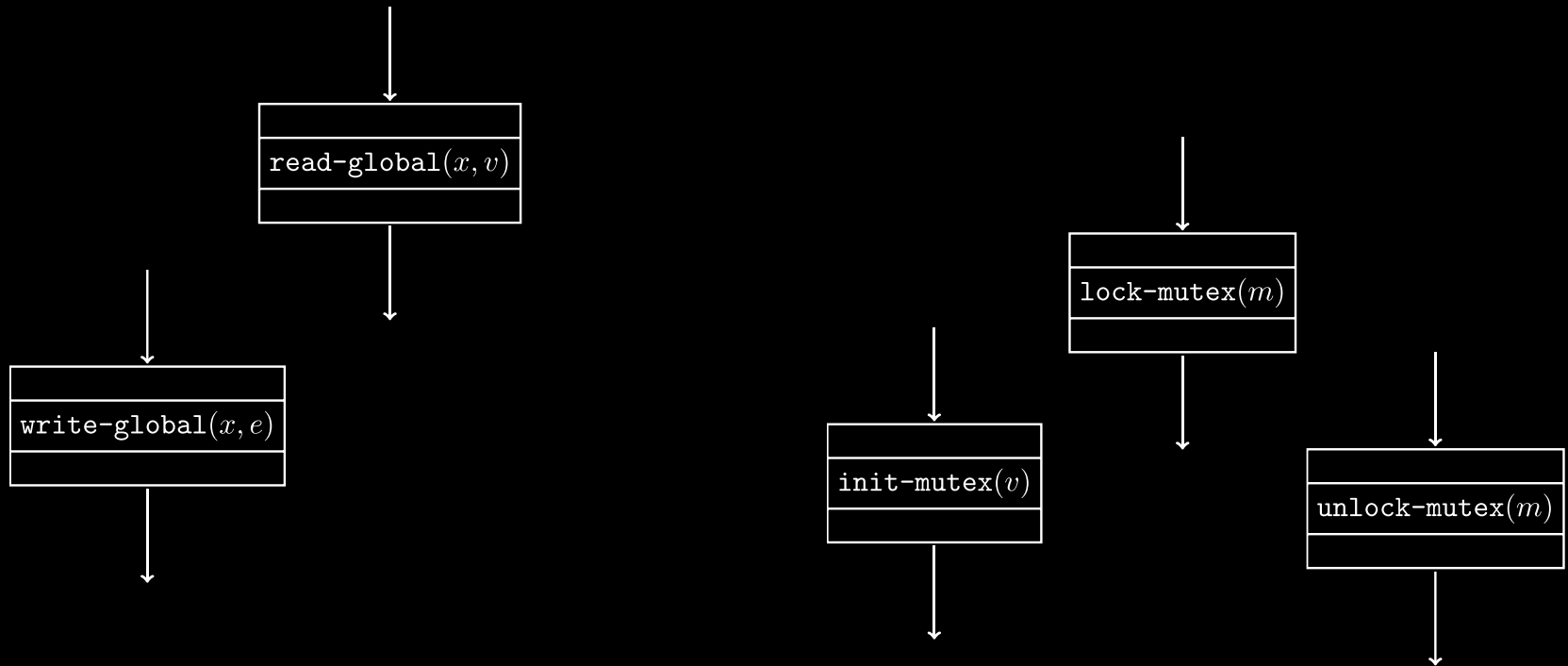
Command Transitions: Agent Management



Command Transitions: Message Passing



Command Transitions: Shared Memory

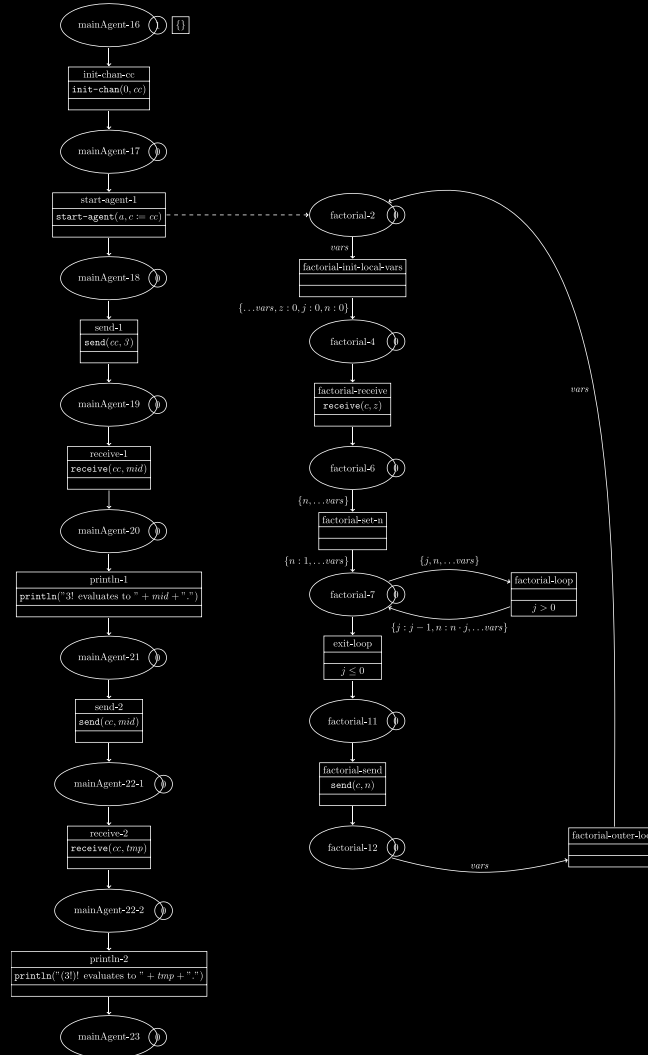


CPPN Syntax restrictions

All agent management must be done through the appropriate command transitions:

- Petri transitions must have exactly one incoming and outgoing arc
- initial marking must have exactly one token

Example



```

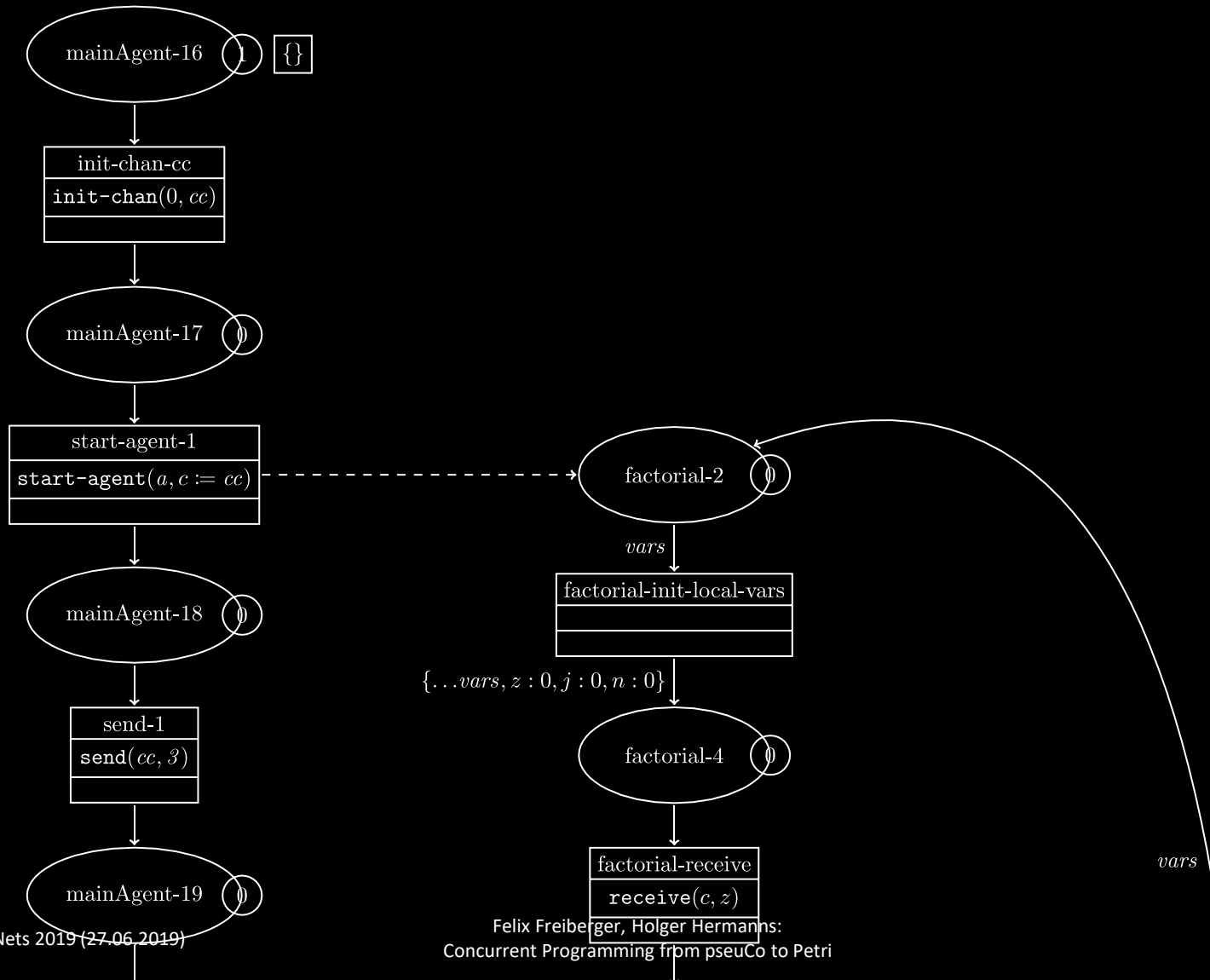
pseuCo by Example

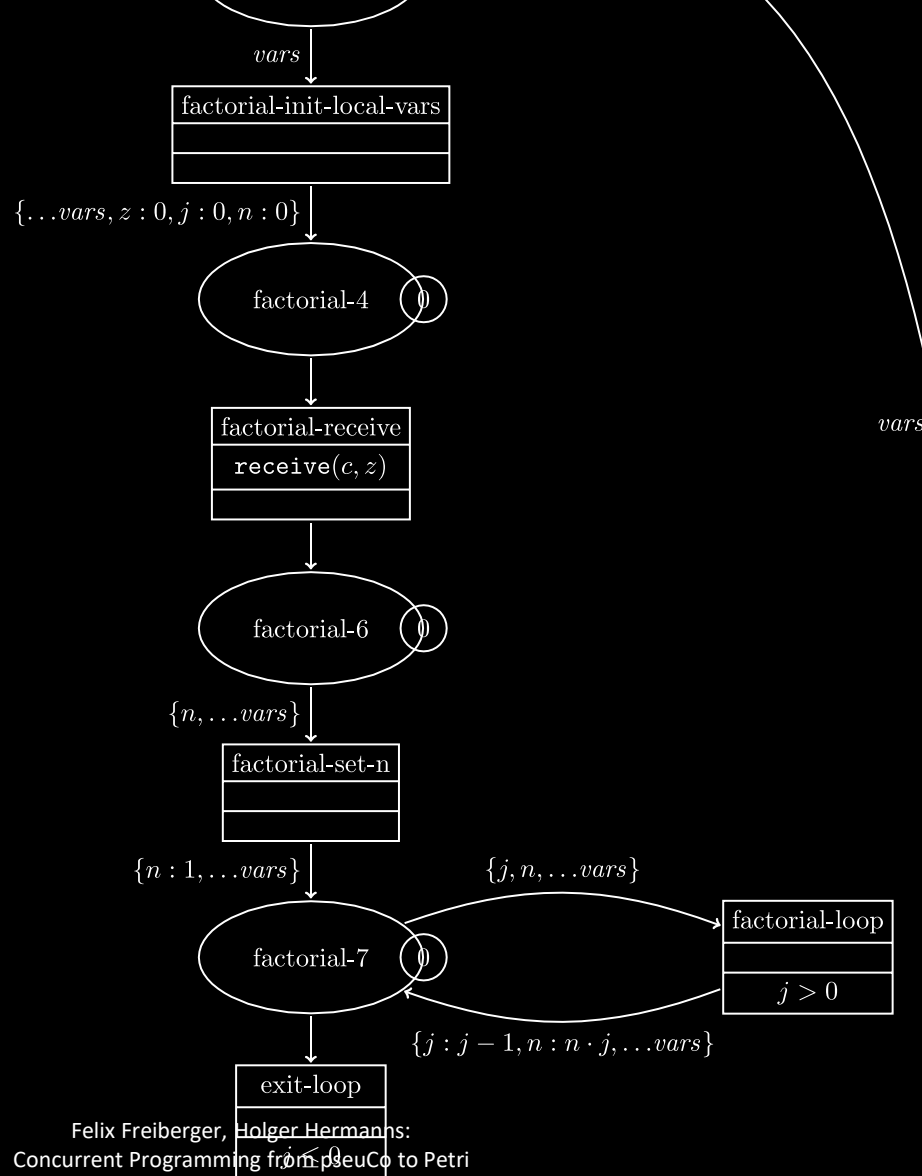
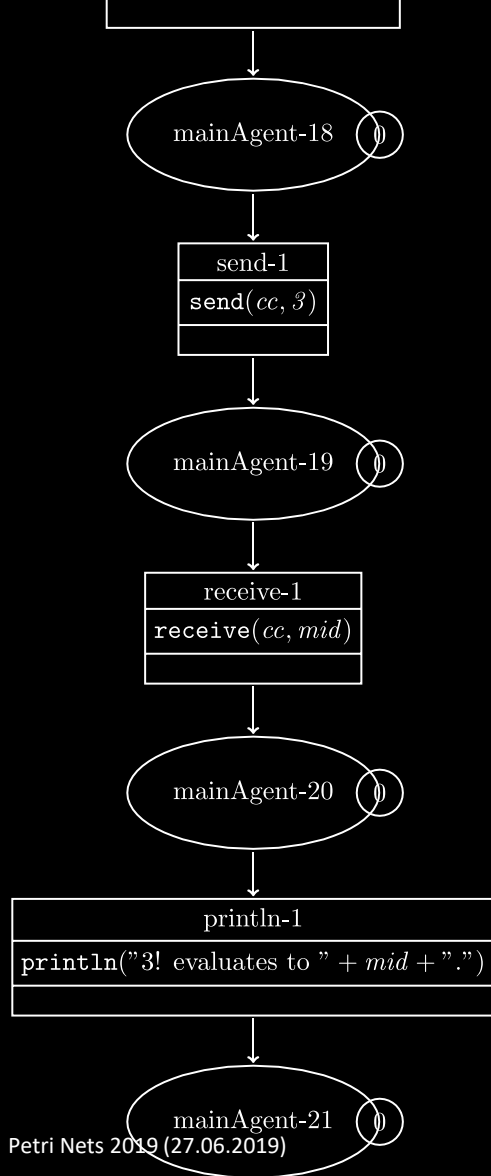
void factorial(intchan c) {
  int z, j, n;
  while (true) {
    z = <? c; // receive input
    n = 1;
    for (j = z; j > 0; j--) {
      n = n*j;
    }
    c <! n; // send result
  }
};

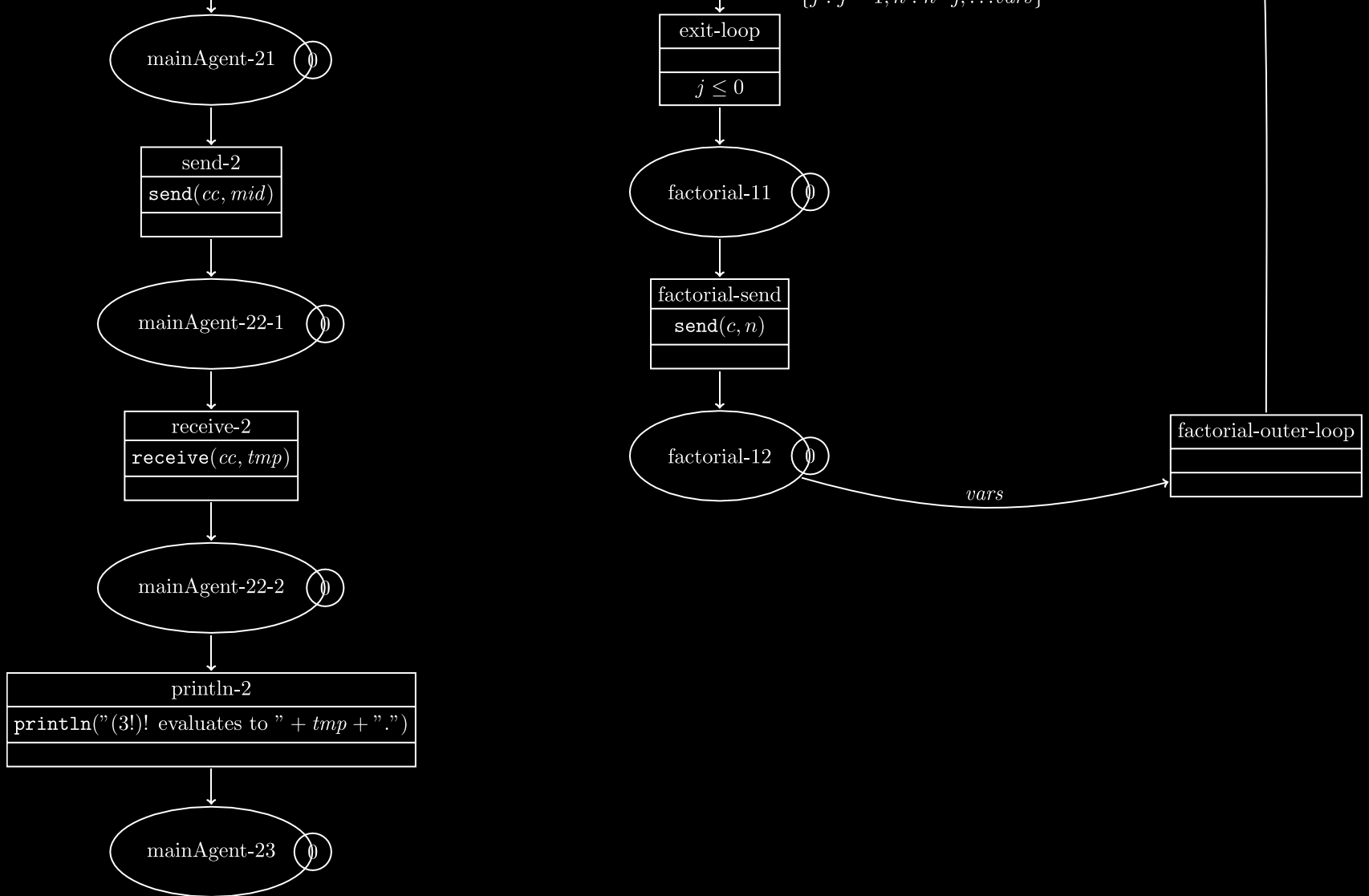
mainAgent {
  intchan cc;
  agent a = start(factorial(cc));
  cc <! 3;
  int mid = <? cc;
  println("3! evaluates to " + mid + ".");
  cc <! mid;
  println("3! evaluates to " + (<? cc) + ".");
};
}

Petri Nets 2019 (27.06.2019) Felix Freiberger, Holger Hermanns Concurrent Programming from pseuCo to Petri Slide 9

```








Semantics

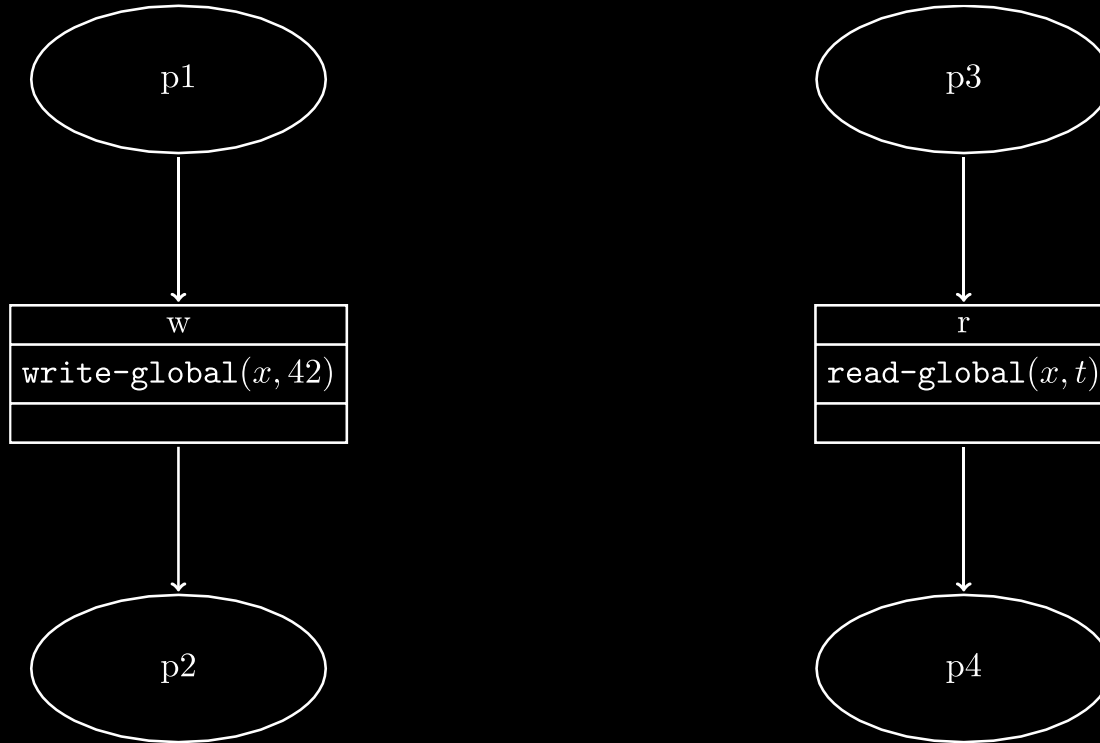
CPPNs unfold to regular CPNs.

This unfolding...

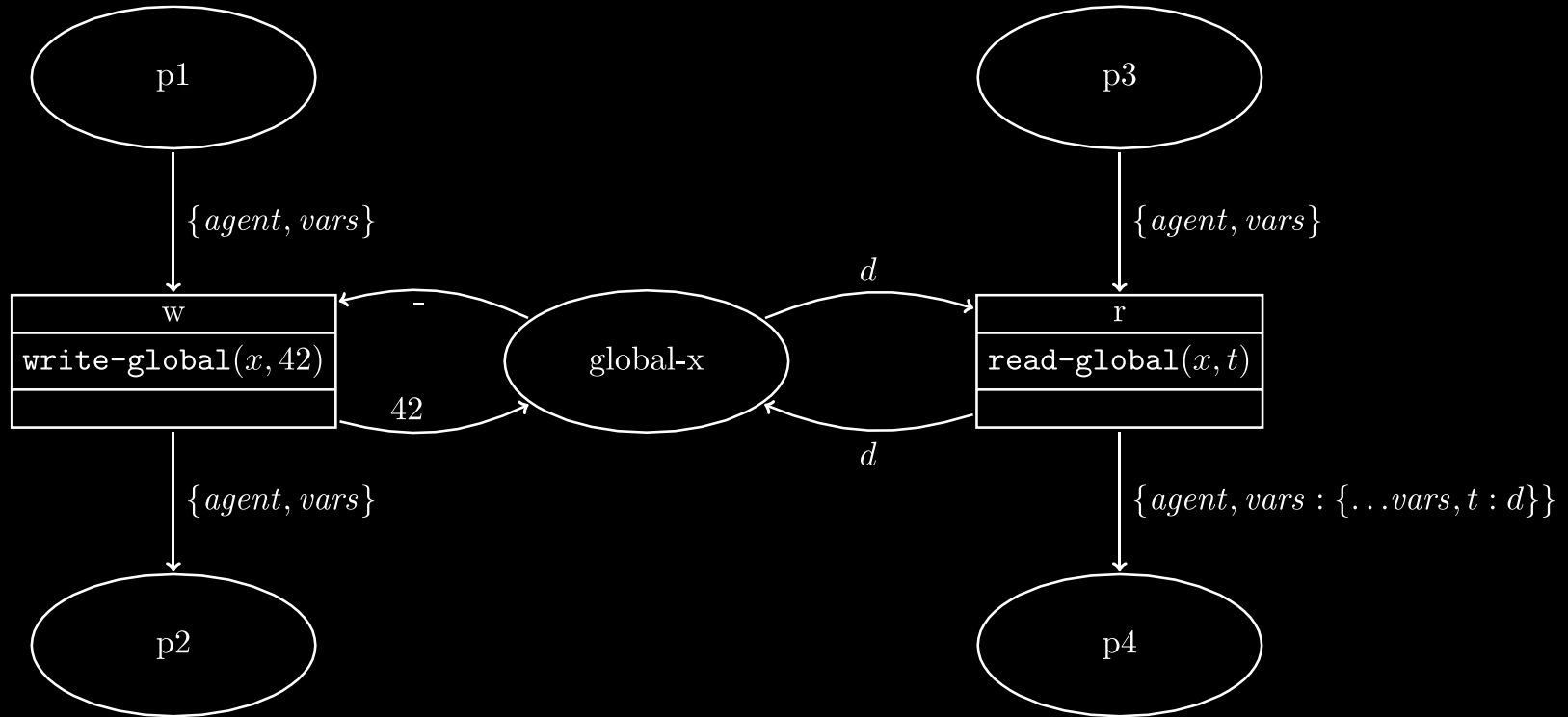
- adds agent identifiers to each token
 - replaces most command transitions with fixed static constructs
 - inserts transitions linking `call` \leftrightarrow `return`
and `send` \leftrightarrow `receive`
- 

Complexity: $\mathcal{O}(n^2)$ in number of these transitions,
 $\mathcal{O}(n)$ for everything else

Unfolding Example

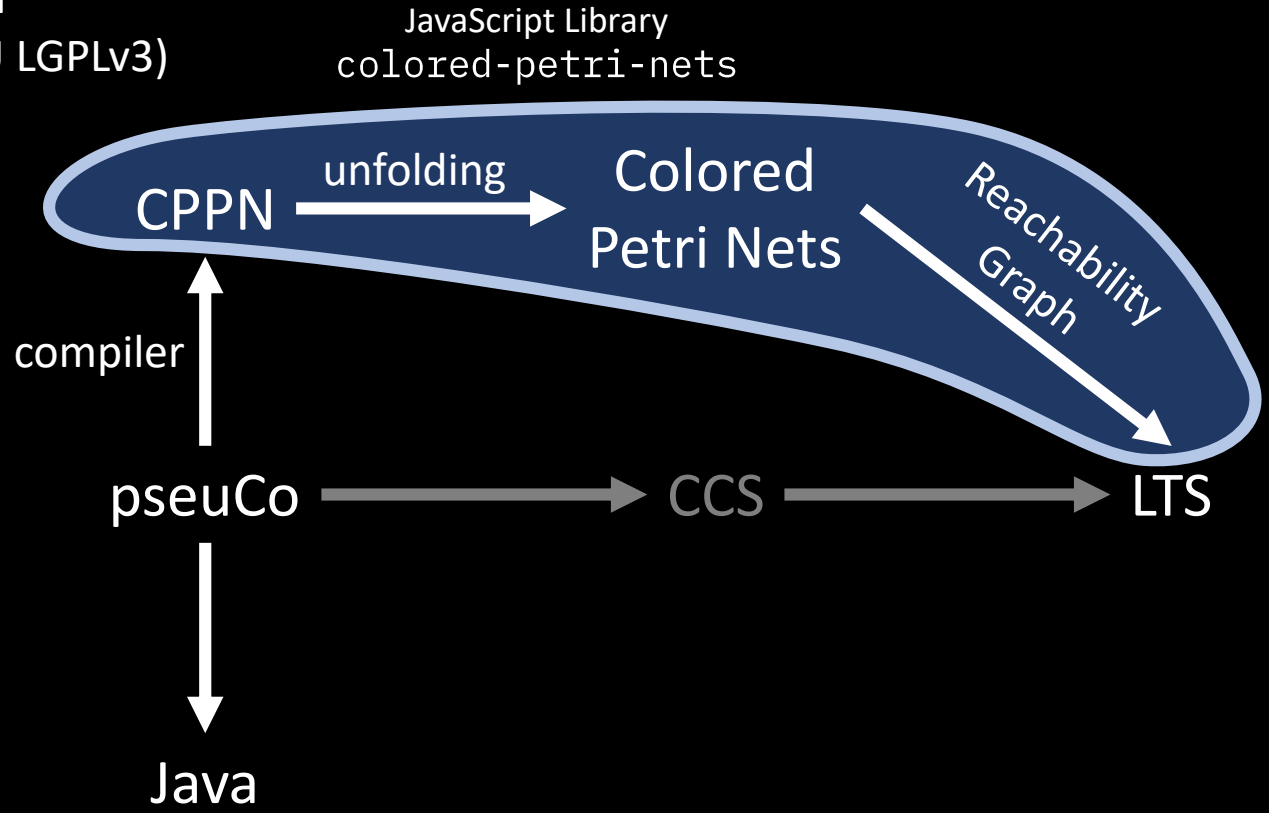


Unfolding Example



Implementation

- published on NPM
- open source (GNU LGPLv3)

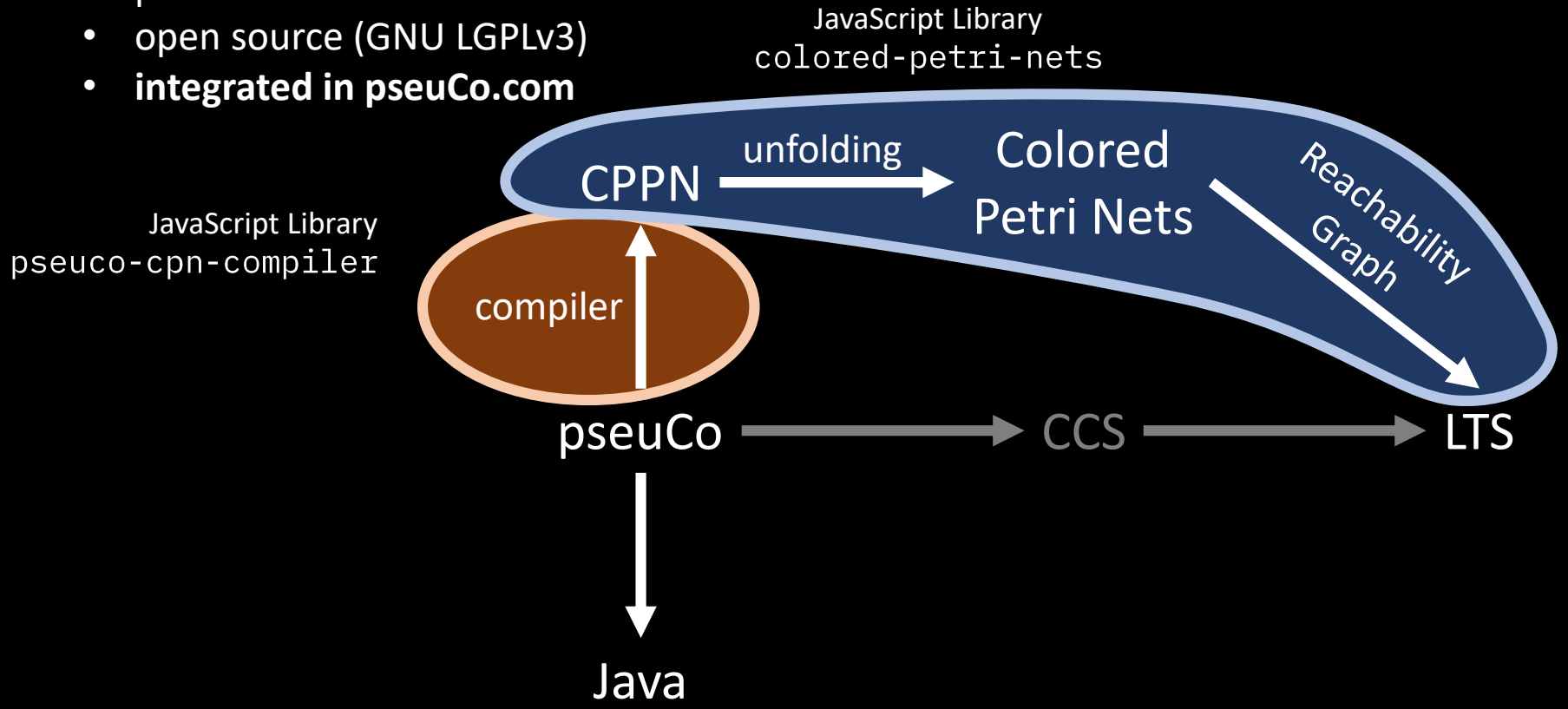


Compiling pseuCo to CPPN

- straightforward compiler
- translates AST nodes into CPPN transitions
- about 2.500 lines of code, with full support for
 - channels
 - select-case
 - global variables
 - locks
 - arrays **new**
 - structures **new**
 - monitors **new**
 - condition synchronization **new**

Implementation

- published on NPM
- open source (GNU LGPLv3)
- **integrated in pseuCo.com**



Editing

First Message Passing

pseuCo Duplicate

pseuCo → CCS → LTS

pseuCo

```

1 void factorial(intchan c) {
2   ...int z, j, n;
3   ...while (true) {
4     .....z = <? c; // receive input
5     .....
6     .....n = 1;
7     .....for (j = z; j > 0; j--) {
8       .....n = n*j;
9     .....}
10    .....
11    .....c <! n; // send result
12    .....};
13  }
14
15  mainAgent {
16    ...intchan cc;
17    ...agent a = start(factorial(cc));
18    ...cc <! 3;
19    ...int mid = <? cc;
20    ...println("3! evaluates to " + mid + ".");
21    ...cc <! mid;
22    ...println("(3)! evaluates to " + (<?
23    cc) + ".");
24  }

```

No issues.

CCS

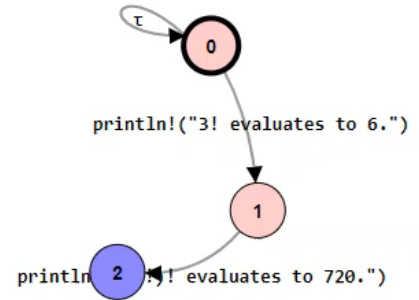
```

1 Channel_cons[i] := channel_create!(i).
  (Channel_cons[i-1])
2 AgentManager[next_i] := agent_new!(next_i).
  (AgentManager[next_i+1])
3 Agent_factorial := agent_new?i.
  (start_factorial!(i).(start_set_arg(i)?
  starter.(start_set_arg(i)?a1.
  (Agent_factorial |
  -(Proc_factorial[i, a1]; agent_terminate(i)!.
  (0))))))
4
5 Proc_factorial[a, - $c] := Proc_factorial_1[a, - $
  c, 0, 0, 0]
6 Proc_factorial_1[a, $c, -$z, -$j, -$n] := when (t
  rue) (τ.(receive($c)?$0.
  (Proc_factorial_3[a, -$c, $0, $0, -1])) + when
  (!true) (Proc_factorial_6[a, $c, -$z, -$j, -$n])
7 Proc_factorial_2[a, $c, -$z, -$j, -$n] := Proc_fa
  ctorial_3[a, -$c, -$z, -$j-1, -$n]
8 Proc_factorial_3[a, $c, -$z, -$j, -$n] := when (!
  ($j>0)) (Proc_factorial_4[a, $c, -$z, -$j, -$n])
  + when ($j>0) (τ.
  (Proc_factorial_2[a, -$c, -$z, -$j, -$n*$j]))
9 Proc_factorial_4[a, $c, -$z, -$j, -$n] := when ($
  c>=0) (put($c)!($n).
  (Proc_factorial_5[a, -$c, -$z, -$j, -$n])) + when
  ($c<0) (receive($c)!($n).
  (Proc_factorial_5[a, -$c, -$z, -$j, -$n]))
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  ctorial_1[a, -$c, -$z, -$j, -$n]

```

No issues.

LTS



Return Collapse all Expand all

Petri Nets to the Rescue

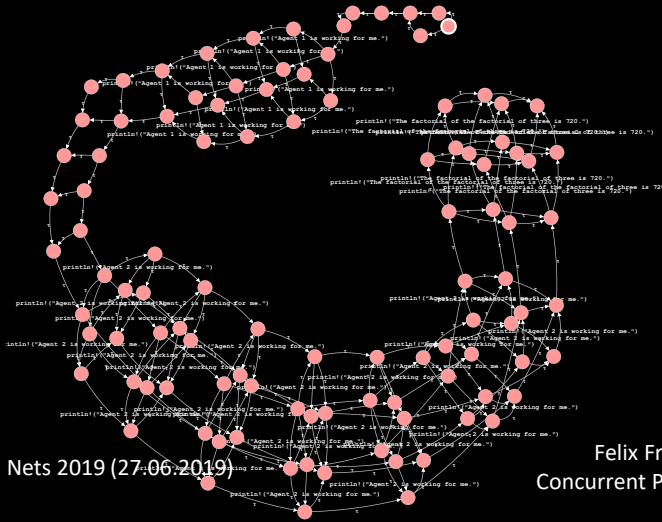
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- hard to understand ✓?
 - control flow hard to see (goto-style spaghetti code) ✓
 - helper constructs (agent management, channels, arrays, ...) ✓
 - low-level hacks visible in the code (e.g. for channels) ✓
- no proper debugging support in pseuCo.com ←
- no true concurrency due to CCS interleaving semantics ✓
- ...and lack of Petri Nets! ✓

Debugging pseuCo: Previous Options

Debugging with CCS and LTS

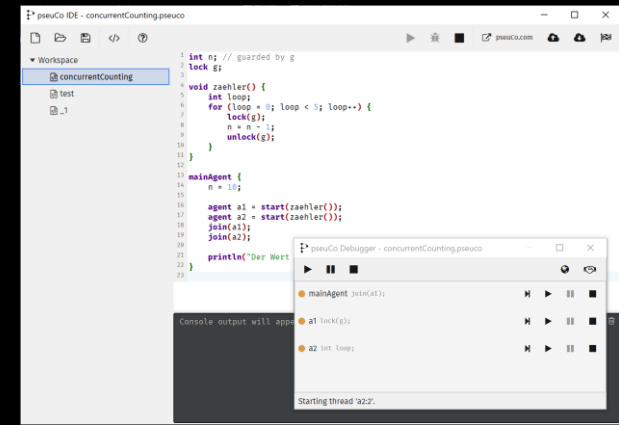
- nondeterminism fully preserved
- **state identifiers are CCS terms, very hard to understand program state**



Petri Nets 2019 (27.06.2019)

Debugging with pseuCo IDE

- traditional, IDE-style debugging
- based on pseuCo → Java semantics
- **no control over nondeterminism**

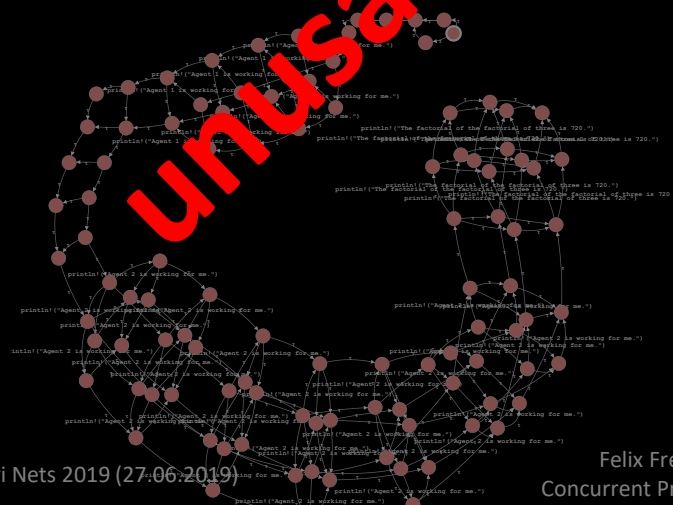


Felix Freiberger, Holger Hermanns:
Concurrent Programming from pseuCo to Petri

Debugging pseuCo: Previous Options

Debugging with CCS and LTS

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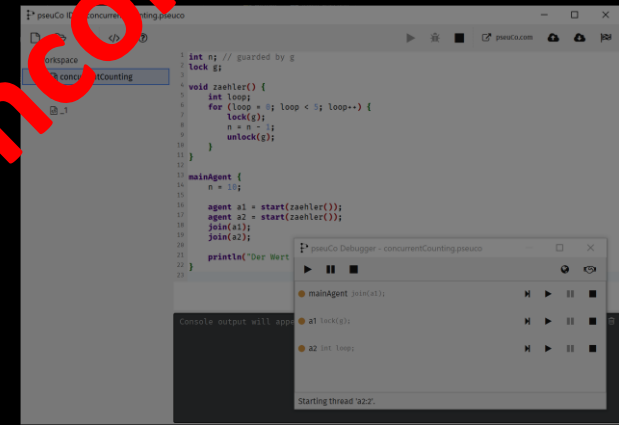


Debugging with pseuCo IDE

- traditional, IDE-style debugging
- based on pseuCo → Java semantics
- **no control over nondeterminism**

unusable

incomplete



Debugging pseuCo: Using Petri Nets

- in principle, the reachability graph is perfect for debugging
 - nondeterminism fully preserved
 - markings are easier to comprehend than CCS terms (once you have understood the Petri net)
- but it requires understanding of Petri nets and compiler internals

**Can we build a debugger that feels like an IDE,
but is actually just walking the reachability graph?**

Editing

First Message Passing

pseuCo Duplicate

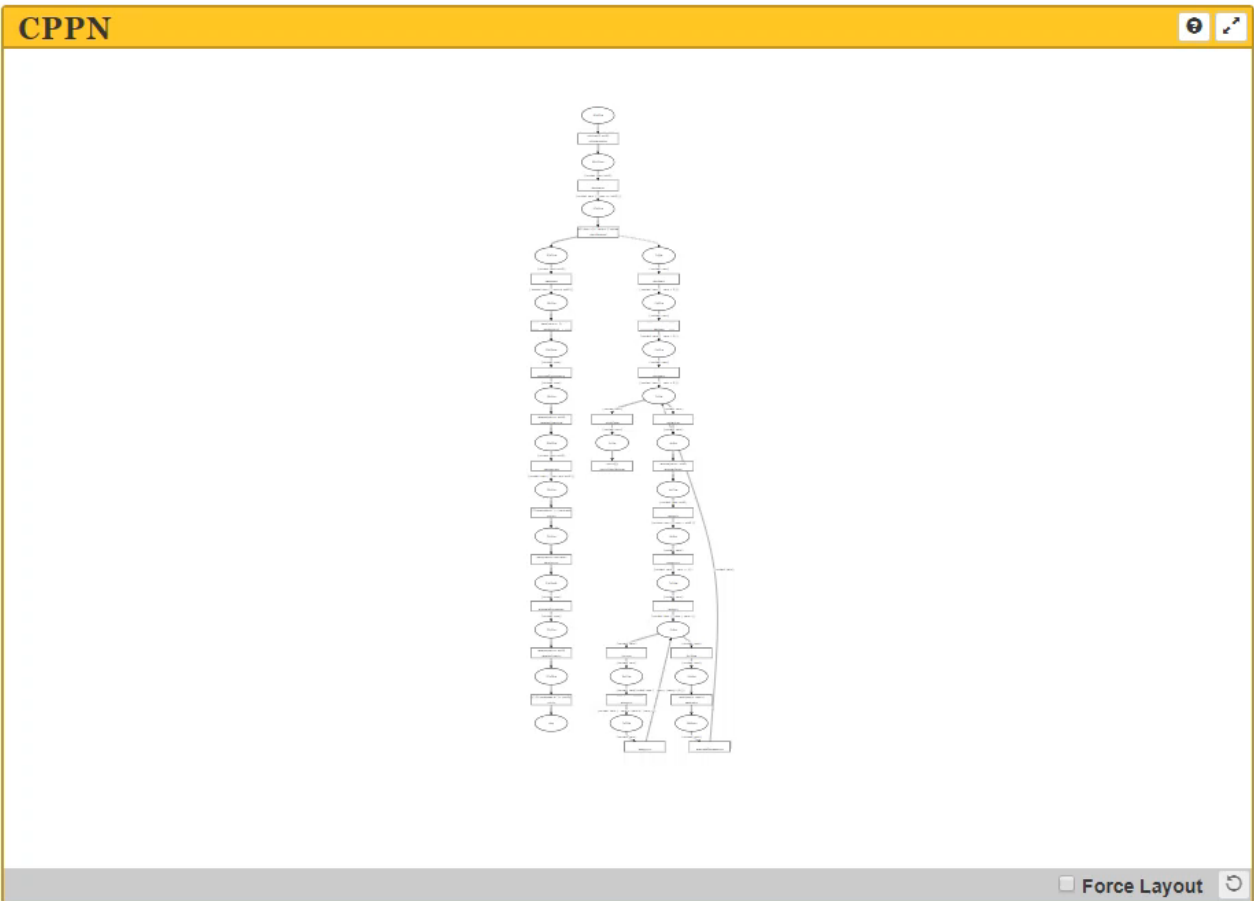
pseuCo -> CPPN -> CPN -> LTS (experimental) pseuCo Debugger (experimental)

```

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3   ...while (true){
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19    ...int mid = <? cc;
20    ...println("3!");
21    ...evaluates to " + mid + ".";
22    ...cc <! mid;
23    ...println("(3)!");
24    ...evaluates to " + (<? cc) + ".";
25  }

```

No issues.

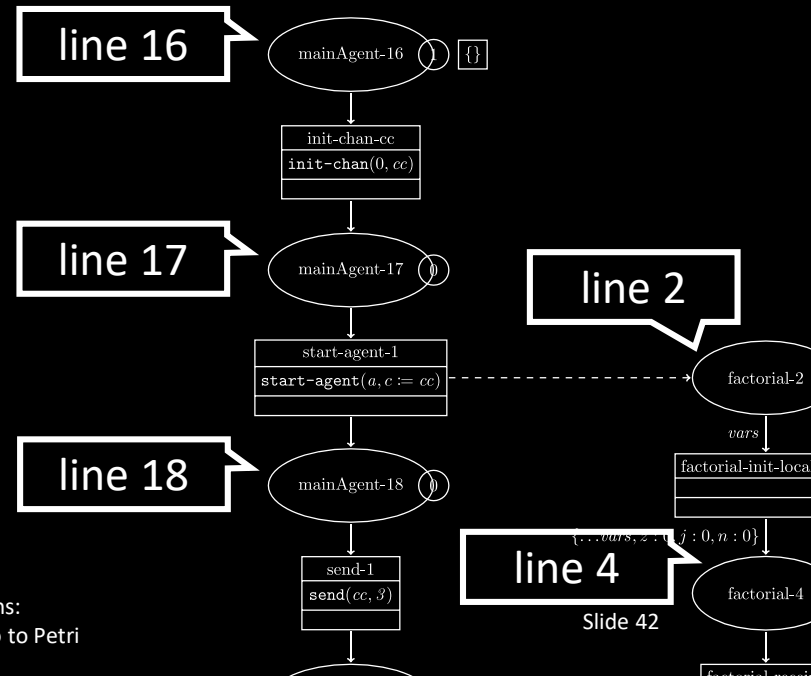


Force Layout

CPN LTS

pseuCo Debugger Internals

- walks reachability graph
- net annotated with auxiliary information (globally and per place)



Summary

- CPPN: a CPN-like formalism for concurrent programs, designed for teaching
- unfolding to regular CPNs
- pseuCo \rightarrow CPPN
- implemented and integrated into pseuCo.com
- Debugging support in pseuCo.com

More information
at the tool exhibition!